## **PERSONAL - Finding Loop Invariants**

## **Tips for finding Loop Invariants**

## **TL;DR: Short Summary**

- Tip 1: invariant must contain variable we operate on and must be as precise as possible
- Tip 2: "loop-carried" variables must be included in the loop invariant
- Tip 3: create generalized tables to visualize the connections between variables inside the loop
- Tip 4: there must be a relation between the variables needed to calculate *x* in the loop and after the loop
- **Tip 5**: if the loop condition contains **inequality**, the counter should be "**limited** on the opposite side" in order to reach **equality**
- **Tip 6**: if certain program inputs are **restricted**, these restrictions about input variables should be included in the loop invariant
- **Tip 7**: if the variable in our loop invariant **depends on some other value**, it needs to be included in some way in our loop invariant
- Tip 8: when proving termination, r must be included in the loop invariant
- Tip 8.5: when proving termination, all variables required for calculating *r* must be included in the loop invariant

## **Detailed Tips**

- Tip 1: invariant must contain variable we operate on and must be as precise as possible
  - example:
    - $i \geq 0$  would be a **bad invariant**, since we cannot imply x = 2n
    - similarly,  $x \ge 0$  would also be a **bad invariant**, since we still cannot imply x = 2n
    - x = 2i, however, would suffice



- **Tip 2**: **"loop-carried" variables** (*variables which depend on their value from the previous loop iteration*) must be included in the loop invariant
  - example:
    - $x = \sum_{k=0}^{i} 3k$  would be a **bad invariant**, since we know nothing about y
    - $x = \sum_{k=0}^{i} 3k \wedge y = 3i$  would **work**, however



• if, instead, y was not dependent on its previous value,  $x = \sum_{k=0}^i 3k$  would suffice



- Tip 3: create generalized tables to visualize the connections between variables inside the loop
  - example:

$ \sum_{k=0}^{\infty} k $					
#	0	1	2		i
i	0	1	2		i
х	2	2	2		2
		+ 2 * 1	+ 2 * 1		+ 2 * 1
			+ 2 * 1 * 2		+ 2 * 1 * 2
					+
					+2*1**i
У	1	1	1 * 2		1 * 2 * * i

• we see that y = i! and  $x = 2\sum_{k=0}^{i} k!$ 

as such, a good invariant would be the conjunction of these two, namely  $x = 2\sum_{k=0}^{i} k! \wedge y = i!$ 



- Tip 4: there must be a relation between the variables needed to calculate x in the loop and after the loop
  - example:
    - $x = 2i^2 32$  cannot imply  $2n^2 + 16|n|$ ; there needs to be some **connection** between n and i or k
    - we see that k=|n|+4, and incorporating this into our loop invariant would allow us to make the needed connection ( $I\equiv k=|n|+4\wedge x=2i^2-32$ )



Tip 5: if the loop condition contains inequality, the counter should be "limited on the opposite side" in order to reach equality (e.g. if *i* is decremented, include *i* ≥ 0, else if *i* is incremented, include *i* ≤ *n*)

```
• example:
```

- compared to previous examples, x = 2i would **not** suffice, since we cannot imply that  $i \ge n$  (we need to show that, at that moment, i = n holds)
- we make our assertion stronger:  $(i < n \land x = 2i) \lor (i \ge n \land x = 2n) \Longleftarrow (i < n \land x = 2i) \lor (i = n \land x = 2n)$
- eventually, we reach  $x = 2i \land i \le n$ , but x = 2i still does not imply this assertion, meaning we must include  $i \le n$  in our assertion ( $I \equiv x = 2i \land i \le n$ )



- **Tip 6**: if certain program inputs are **restricted** (e.g. having to take the absolute value of an input *n* to only allow positive integers), these restrictions about input variables should be included in the loop invariant
  - example:
    - this program doesn't work for some negative *n*...



 as such, some construct is most likely included before the loop to handle problematic inputs, which **must be accounted for** in the loop invariant



- Tip 7: if the variable in our loop invariant depends on some other value, it needs to be included in some way in our loop invariant if the value constantly fluctuates between one value and another, a simple formula will do, otherwise, do case-by-case if-then-else analysis (e.g. (c1 ⇒ ...) ∧ ... ∧ (c<sub>n</sub> ⇒ ...))
  - example:
    - x is dependent on the value b, which is either 0 or 1 in any iteration (but 0 when we're done)
    - as such, a simple formula  $I\equiv x=4i+b$  suffices



- **Tip 8**: when proving termination, *r* **must be included** in the loop invariant
- **Tip 8.5**: when proving termination, **all variables required** for calculating *r* must be included in the loop invariant